Reliable Data Transmission and Energy Efficient Routing Protocol for WSNs
Shailendra\textsuperscript{1}, Munna Singh Kushwaha\textsuperscript{2}

\textsuperscript{1}M.Tech (ECE), \textsuperscript{2}Head of Department (ECE)
Somany (PG) Institute of Technology & Management
Rewari, Haryana, India

Abstract: In WSN, hierarchical routing architecture partitions the whole network into a group of cluster and only cluster head is responsible for forwarding the data to base station or sink directly. In these networks some of the nodes become cluster heads, aggregate the data of their cluster members and transmit it to the sink. We assume that a percentage of the population of sensor nodes is equipped with additional energy resources—this is a source of heterogeneity which may result from the initial setting or as the operation of the network evolves. We also assume that the sensors are randomly distributed and are not mobile, the coordinates of the sink and the dimensions of the sensor field are known. We show that the behavior of such sensor networks becomes very unstable once the first node dies, especially in the presence of node heterogeneity. In this work, we have considered the concept of heterogeneity as in Stable Election Protocol abbreviated as SEP. When a percentage of the population of sensor nodes is equipped with additional energy resources, or extra energy is provided to some nodes, then the network is known as heterogeneous network. We also assume that the sensor nodes are randomly distributed in a region and are not mobile, the coordinates of the base station and the dimensions of the sensor field are known. In this work, we have designed a routing protocol based on SEP and make it reliable as well as energy efficient by considering the case of residual energy as well as sleep/awake scheduling and results are compared with existing routing protocol of same category such as SEP.

Keywords: Heterogeneous wireless sensor network, SEP, energy efficient, reliable

1. INTRODUCTION
With a sensor network, a user should be able to task some sensors to monitor specific events, and know when interested events happen in the interested field [1]. Thus, the sensor network builds a bridge between the real world and computation world. Each node typically consists of the five components: sensor unit, analog digital convertor (ADC), central processing unit (CPU), power unit, and communication unit [1,2]. The sensor unit is responsible for collecting information as the ADC requests, and returning the analog data it sensed. ADC is a translator that tells the CPU what the sensor unit has sensed, and also informs the unit what to do. Communication unit is tasked to receive command or query from, and transmit the data from CPU to the outside world. CPU is the most complex unit [3]. It interprets the command or query to ADC, monitors and controls power if necessary, processes received data, computes the next hop to the sink, etc. Many Routing protocols are existent in the wireless sensor network [4].

2. ROUTING TECHNIQUES IN WSNs
Depending on how the sender of a message gains a route to the receiver, routing protocols can be classified into three categories. Flat routing architecture, hierarchical routing architecture, and location based architecture are one of those.

(a) Flat routing architecture: The first category of routing protocols is the multihop flat routing protocols. In flat routing, each node typically plays the same role and sensor nodes collaborate to perform the sensing task. Due to the large number of such nodes, it is not feasible to assign a global identifier to each node. This consideration has led to data-centric routing, where the BS sends queries to certain regions and waits for data from the sensors located in the selected regions. Since data is being requested through queries, attribute-based naming is necessary to specify the properties of data [1-5].

(b) Hierarchical routing architecture: Hierarchical or cluster-based routing methods, originally proposed in wire line networks, are well-known techniques with special advantages related to scalability and efficient communication. As such, the concept of hierarchical routing is also utilized to perform energy efficient routing in WSNs. The main idea in this class of protocols is that every sensor node within a WSN is grouped along with some other of its neighboring nodes so that to constitute a specific cluster. Data collected by the sensor nodes belonging to a cluster are not directly transmitted to the base station. Instead, a node of the cluster, called cluster head, collects these data and forwards them to the base station after possibly having performed appropriate data aggregation [6-8]. In this way, the number of transmitted messages to the base station is reduced and considerable power conservation is achieved. Hierarchical routing can be further two types as Centralized and non-centralized hierarchical routing. In non-centralized type of each node self configures for the cluster head. Whereas, in centralized type of hierarchical routing base station takes this responsibility.

(c) Location based routing architecture: In this kind of routing, sensor nodes are addressed by means of their locations. The distance between neighboring nodes can be estimated on the basis of incoming signal strengths. Relative coordinates of neighboring nodes can be obtained by exchanging such information between neighbors [9,10].
3. PROPOSED ALGORITHM

The foundation of proposed protocol lies in the realization that the base station is a high-energy node with a large amount of energy supply. Thus, proposed protocol utilizes the base station to control the coordinated sensing task performed by the sensor nodes.

In proposed protocol, the following assumption are to be considered:

- A fixed base station is located in the middle of the region.
- The nodes are equipped with power control capabilities to vary their transmitted power.
- Each node senses the environment at a fixed rate and always has data to send to the base station.
- All sensor nodes are immobile.

1. The radio channel is supposed to be symmetrical. Thus, the energy required to transmit a message from a source node to a destination node is the same as the energy required to transmit the same message from the destination node back to the source node for a given SNR (Signal to Noise Ratio). Moreover, it is assumed that the communication environment is contention and error free. Hence, there is no need for retransmission.

2. A percentage of the population of sensor nodes is equipped with more energy resources than the rest of the nodes. Let \( m \) be the fraction of the total number of nodes \( n \), which are equipped with \( \alpha \) times more energy than the others. We refer to these powerful nodes as advanced nodes, and the rest \( (1-m)n \) as normal nodes. We assume that all nodes are distributed randomly over the sensor field.

3. Suppose that \( Eo \) is the initial energy of each normal sensor. The energy of each advanced node is then \( Eo^\alpha \). The total (initial) energy of the new heterogeneous setting is equal to: \( n \times Eo^\alpha \). So, the total energy of the system is increased by a factor of \( 1+\alpha \).

\( i \) each normal node becomes a cluster head once every \( 1/pnrm \) rounds per epoch; 
\( ii \) each advanced node becomes a cluster head exactly once every \( 1/padv \) rounds per epoch; 
\( iii \) the average number of cluster heads per round per epoch is equal to \( n \times 1/pnrm \)

4. Cluster Head Election for normal nodes is based on following equation:

\[
T(s_{nrm}) = \begin{cases} 
\frac{Eo}{1-pnrm \cdot \frac{1}{1-pnrm}} & \text{if } s_{nrm} \in G' \\
0 & \text{otherwise}
\end{cases}
\] (1)

5. Cluster Head Election for advanced nodes is based on following equation:

\[
T(s_{adv}) = \begin{cases} 
\frac{Eo}{1-padv \cdot \frac{1}{1-padv}} & \text{if } s_{adv} \in G'' \\
0 & \text{otherwise}
\end{cases}
\] (2)

where \( G'' \) is the set of advancednodes that have not become cluster heads within the last \( 1/padv \) rounds of the epoch, and \( T(s_{adv}) \) is the threshold applied to a population of \( n \times m \) (advanced) nodes. This guarantees that each advanced node will become a cluster head exactly once every \( 1/padv \cdot ((1+\alpha m)/((1+\alpha)) \) rounds.

6. Based on above equations and conditions, nodes sends the data to their respective cluster heads and energy consumption will be calculated.

7. Cluster Head will aggregate the data and send it to the base station and energy consumption will be calculated for each node and cluster heads.

8. In round 2, the nodes will become cluster heads according to probability condition i.e. according to minimum distance from base station and threshold energy.

9. After selection of cluster heads, Nodes sends the data to their respective cluster heads, that will be selected according to the minimum distance of a particular node from cluster heads and energy consumption will be calculated.

10. Cluster Head will aggregate the data and send it to the base station and energy consumption will be calculated.

11. Ten nodes will also go in sleep mode to enhance the network lifetime if their energy is less than 1 nJ. If the numbers increase then ten, then the nodes will come in active mode and send the data to nearby cluster head.

12. This process will be repeated until the whole network gets down or number of rounds finished. Performance will be evaluated according to parameters like network lifetime, energy dissipation, data packets sent etc.

4. IMPLEMENTATION & RESULTS

4.1 Parameter Value

<table>
<thead>
<tr>
<th>Table 1 Network Parameters</th>
</tr>
</thead>
<tbody>
<tr>
<td>Network field:</td>
</tr>
<tr>
<td>N (Number of nodes):</td>
</tr>
<tr>
<td>Initial energy:</td>
</tr>
<tr>
<td>Eelec (ETx&amp;ERx):</td>
</tr>
<tr>
<td>ε fs (free space):</td>
</tr>
<tr>
<td>εmp :</td>
</tr>
<tr>
<td>EDA :</td>
</tr>
<tr>
<td>Data packet size:</td>
</tr>
<tr>
<td>Tool used</td>
</tr>
</tbody>
</table>

4.2 Results

Refer figure 1, It shows the plotting of 100 nodes that deployed randomly in an area of 200 m X 200 m. ‘o’ symbol denoted as normal nodes, ‘+’ symbol denoted as advanced nodes with extra energy and ‘x’ indicates position of base station.
Figure 1: Plotting of nodes for SEP protocol

Figure 2: Plotting of nodes for proposed protocol

Refer figure 2, It shows the plotting of 100 nodes that deployed randomly in an area of 200 m X 200 m. ‘o’ symbol denoted as normal nodes, ‘+’ symbol denoted as advanced nodes with extra energy and ‘x’ indicates position of base station.

Figure 3: Number of Rounds vs Number of Nodes Alive

Figure 3 shows the comparison of routing protocols Stable Election Protocol (SEP) and proposed protocol in terms of Number of nodes alive. As we can see in figure, when number of rounds increases; the numbers of nodes alive are comparatively less in SEP as compared to proposed protocol. Figure 3 shows the overall lifetime of the network. Here, we can observe that proposed routing technique performs better as compared to SEP.

Figure 4: Number of Rounds vs Number of Nodes Dead

Figure 4 shows the comparison of routing protocols Stable Election Protocol (SEP) and proposed protocol in terms of Number of nodes dead. As we can see in figure, when number of rounds increases; the numbers of nodes dead are comparatively less in proposed protocol as compared to SEP. Figure 4 shows the overall lifetime of the network. Here, we can observe that proposed routing technique performs better as compared to SEP.

Figure 5: Number of Rounds vs Data Packets sent to base station (SINK)

Figure 5 shows how much data will be sent from nodes to SINK. From figure 5, we can observed that, in SEP protocol data sent to base station is relatively less as compared to proposed routing technique proposed protocol.

Figure 6: Number of Rounds vs Energy Consumption in each round
Figure 6 shows the lifetime of the network. It shows that how energy of the network consumes round by round and finally whole network goes down. It can be observed from the figure 6 that, proposed protocol consumes less energy and sustain more number of rounds as compare to SEP protocol.

5. CONCLUSION AND FUTURE WORK
The proposed routing protocol which is hierarchical routing based with the whole control to the base station or we can say that base assisted. In proposed routing technique, the base station first collects information about the logical structure of the network and residual energy of each node. So, with the global information about the network base station does cluster formation better in the sense that it has information about the residual energy of each node. In proposed protocol, the concept of sleep nodes also enhances the network lifetime. Finally, proposed routing technique is compared with already developed routing protocol Stable Election Protocol (SEP) by the help of MATLAB. A comparison between two is done on the basis of data packet sent and the system lifetime of network.

In WSN, hundreds or thousands of sensor nodes are randomly scattered in the sensor field. These nodes sense the data and send this sensed data to the cluster head (in case of hierarchical routing) or directly to the base station according to the TDMA (time division multiplexing access) given by cluster head or base station respectively. But there is no security and authentication while communicating. So this can be another research area where this can be considered. So in future, security can be applied to proposed routing technique.

REFERENCES